

COMPLEXITY OF LATTICE—CONFIGURATIONS

by

T. G. TARJÁN

Introduction

Let $Y = \{y_1, y_2, \dots, y_r\}$ and $Z = \{z_1, z_2, \dots, z_s\}$ be finite sets. Let us define the set \mathcal{R} of the discrete rectangles:

$$\mathcal{R} = \{R: R = U \times V; U, V \neq \emptyset; U \subset Y; V \subset Z\}$$

where $U \times V$ denotes the Cartesian product of U and V : $U \times V = \{(y, z): y \in U, z \in V\}$. For the power set $\mathcal{P} = 2^{(Y \times Z)} = \{P: P \subset Y \times Z\}$ the relation $\mathcal{R} \subset \mathcal{P}$ holds. Let us denote the complexity of the rectangle $R = U \times V \in \mathcal{R}$ by $\pi(R)$, which is defined by $\pi(R) = |U| + |V|$.

Let us extend the complexity function π for \mathcal{P} in the following way: $\pi(P) = \min \left\{ \sum_{R \in \mathcal{Q}} \pi(R): \mathcal{Q} \subset \mathcal{R}, \bigcup_{R \in \mathcal{Q}} R = P \right\}$ if $P \in \mathcal{P}$. Further definitions: $\pi(r, s) = \max_{P \in \mathcal{P}} \pi(P)$, $\pi(r) = \pi(r, r)$.

These concepts arised in the switching theory (see e.g. [7]), but they have some interest in themselves, too.

The problem is to determine the values of $\pi(P)$ for certain sets $P \in \mathcal{P}$ and the value of $\pi(r, s)$.

Hencefort I shall try to determine the complexity for three types of sets P and I shall give an asymptotic lower bound for $\pi(r)$.

To the solution of the first problem I shall generalize a theorem of KATONA [2], in which he has proved a conjecture of EHRENFUCHT and MYCIELSKI. Moreover, I draw two corollaries from this generalization, which do not follow from KATONA's theorem. In this case $\pi(P_1) \sim r \cdot \log_2 r$ holds.

In solving the second problem I use the idea of KATONA and SZEMERÉDI from the proof of Th. 1 in [3] and I get that $\pi(P_2) \sim r \cdot \log_2 r$.

The third problem is the complexity of the Hadamard-matrices [1]. This remained an open question. I shall give only the following lower bound: $\pi(P_3) \cong \equiv (r+1)\sqrt{r}$.

The forth problem is the computing of the value of $\pi(r)$. This is also an unsolved problem. Using LUPANOV's Th. 4 in [5] I shall give only the following asymptotic lower bound for $\pi(r)$ without proof: $\pi(r) \gtrsim \frac{r^2}{\log_2 r \cdot \log_2 \log_2 r}$. Accordingly the complexity of P_1 and P_2 can not be maximal.

1. Complexity of diagonalless squares

Assume $Y=Z$ and consider the set $P_1 = \{p: p = (y_i, y_j), i \neq j, 1 \leq i, j \leq m\}$. What is the value of $\pi(P_1)$?

First I shall prove a lemma which is a generalization of KATONA's theorem in [2] and I shall draw two corollaries of it.

LEMMA 1. Let A_1, A_2, \dots, A_m and B_1, B_2, \dots, B_m be subsets of X ($|X|=n$), such that $A_i \cap B_j \neq \emptyset$ iff $i \neq j$ ($1 \leq i, j \leq m$). Then

$$\sum_{i=1}^m \binom{|A_i| + |B_i|}{|A_i|}^{-1} \leq 1 \quad \text{holds.}$$

PROOF. For arbitrary $i \neq j$, $1 \leq i, j \leq m$ there exist distinct $x_k, x_l \in X$ for which $x_k \in A_i \cap B_j$ and $x_l \in A_j \cap B_i$ hold. Then there is no permutation of the elements $x_1, x_2, \dots, x_n \in X$ for which all the elements of A_i could precede the elements of B_i and all the elements of A_j precede the elements of B_j .

There are $\binom{n}{|A_i| + |B_i|} |A_i|! |B_i|! (n - |A_i| - |B_i|)!$ permutations of the elements $x_1, x_2, \dots, x_n \in X$ in which all the elements of A_i precede the elements of B_i . These permutations are different for different i 's therefore

$$\sum_{i=1}^m \binom{n}{|A_i| + |B_i|} \cdot |A_i|! |B_i|! (n - |A_i| - |B_i|)! \leq n!$$

holds which proves our lemma.

Take the particular case, when $B_i = A_i^c = X \setminus A_i$. In this case the assumptions of the lemma have the form $A_i \not\subset A_j$ ($i \neq j$). That is, Lemma 1 gives the following

Corollary 1. (LUBELL—MESHALKIN—YAMAMOTO inequality [4], [6] and [8]) If A_1, A_2, \dots, A_m are subsets of an n -element set, and $A_i \not\subset A_j$ ($i \neq j$) then

$$\sum_{i=1}^m \binom{n}{|A_i|}^{-1} \leq 1.$$

Before the second corollary I shall prove a simple lemma.

LEMMA 2.

The function $f(x, y) =$

$$= \begin{cases} \frac{1 - \langle x \rangle - \langle y \rangle}{\binom{[x] + [y]}{[x]}} + \frac{\langle y \rangle}{\binom{[x] + \{y\}}{[x]}} + \frac{\langle x \rangle}{\binom{\{x\} + [y]}{\{x\}}} & \text{if } \langle x \rangle + \langle y \rangle \leq 1 \\ & \langle x \rangle \langle y \rangle > 0 \\ \frac{1 - \langle x \rangle - \langle y \rangle}{\binom{\{x\} + \{y\}}{\{x\}}} + \frac{\langle x \rangle}{\binom{[x] + \{y\}}{[x]}} + \frac{\langle y \rangle}{\binom{\{x\} + [y]}{\{x\}}} & \text{if } 0 \leq \langle x \rangle + \langle y \rangle < 1 \end{cases}$$

is convex in the triangle $1 \leq x, 1 \leq y, x + y \leq n$, where

$$[z] = \max_{\substack{a \text{ integer} \\ a \leq z}} a \quad \{z\} = \min_{\substack{a \text{ integer} \\ a \geq z}} a$$

$$\langle z \rangle = z - [z] \quad \langle z \rangle = \{z\} - z$$

Moreover in case of an arbitrary x_0 , $1 \leq x_0 \leq n-1$ $f(x_0, y)$ is a strictly monotonically decreasing function of y in $[1, n-x_0]$. Similarly, $f(x, y_0)$ is also a decreasing function of x in $[1, n-y_0]$, where y_0 is fixed in $[1, n-1]$.

PROOF. Since $f(x, y)$ is a linear and continuous extension of $\binom{x+y}{x}^{-1}$, it is sufficient to show that

- (i) for integer x_0 , $1 \leq x_0 \leq n-1$, $f(x_0, y)$ is a convex function of y , $y \in [1, n-x_0]$
- (ii) for integer y_0 , $1 \leq y_0 \leq n-1$, $f(x, y_0)$ is a convex function of x , $x \in [1, n-y_0]$
- (iii) for integer x_0, y_0 , $2 \leq x_0, 2 \leq y_0, x_0 + y_0 \leq n$,

$$f(x_0, y_0) - f(x_0, y_0 - 1) - f(x_0 - 1, y_0) + f(x_0 - 1, y_0 - 1) \geq 0,$$

because for the plane

$$\begin{aligned} p(x, y) &= f(x_0, y_0) + \frac{f(x_0 + \varepsilon, y_0) - f(x_0, y_0)}{\varepsilon} (x - x_0) + \\ &\quad + \frac{f(x_0, y_0 + \varepsilon) - f(x_0, y_0)}{\varepsilon} (y - y_0) \geq \\ &\geq f(x_0, y_0) + f(x, y_0) - f(x_0, y_0) + f(x_0, y) - f(x_0, y_0) \end{aligned}$$

holds ($\varepsilon = 1$ or -1) from (i) and (ii), and finally from (iii):

$$f(x, y) - p(x, y) \geq f(x, y) - f(x_0, y) - f(x, y_0) + f(x_0, y_0) \geq 0$$

is true and equality holds if (x, y) is equal to (x_0, y_0) , $(x_0 + \varepsilon, y_0)$ or $(x_0, y_0 + \varepsilon)$.

$$\begin{aligned} \text{(i)} \quad &\frac{x_0! y!}{(x_0 + y)!} - \frac{x_0! (y-1)!}{(x_0 + y - 1)!} = \frac{x_0! (y-1)!}{(x_0 + y - 1)!} \left(\frac{y}{x_0 + y} - 1 \right) = -\frac{x_0! (y-1)! x_0}{(x_0 + y)!}, \\ &-\frac{x_0! y! x_0}{(x_0 + y + 1)!} + \frac{x_0! (y-1)! x_0}{(x_0 + y)!} = \frac{x_0! (y-1)! x_0}{(x_0 + y)!} \times \\ &\times \left(-\frac{y}{x_0 + y + 1} + 1 \right) = \frac{(x_0 + 1)! (y-1)!}{(x_0 + y + 1)!} x_0 = x_0 \binom{x_0 + y + 1}{x_0 + 1}^{-1} > 0 \quad \text{if } 2 \leq y. \end{aligned}$$

$$\text{(ii)} \quad \text{Similarly } y_0 \binom{x + y_0 + 1}{y_0 + 1}^{-1} > 0 \quad \text{if } 2 \leq x.$$

$$\begin{aligned} \text{(iii)} \quad &\frac{x_0! y_0!}{(x_0 + y_0)!} - \frac{(x_0 - 1)! y_0!}{(x_0 + y_0 - 1)!} - \frac{x_0! (y_0 - 1)!}{(x_0 + y_0 - 1)!} + \frac{(x_0 - 1)! (y_0 - 1)!}{(x_0 + y_0 - 2)!} = \\ &= \frac{(x_0 - 1)! (y_0 - 1)!}{(x_0 + y_0 - 2)!} \left(\frac{x_0 y_0}{(x_0 + y_0)(x_0 + y_0 - 1)} - \frac{x_0}{x_0 + y_0 - 1} - \frac{y_0}{x_0 + y_0 - 1} + 1 \right) = \\ &= \frac{x_0 y_0 - (x_0 + y_0)^2 + (x_0 + y_0)^2 - (x_0 + y_0)}{\binom{x_0 + y_0 - 2}{x_0 - 1} (x_0 + y_0)(x_0 + y_0 - 1)} = \frac{(x_0 - 1)(y_0 - 1) - 1}{\binom{x_0 + y_0 - 2}{x_0 - 1} (x_0 + y_0)(x_0 + y_0 - 1)} \geq 0, \end{aligned}$$

if $2 \leq x_0$
 $2 \leq y_0$

From the proof of the convexity you can also see the monotony of $f(x, y)$. The proof is completed.

Corollary 2. Let A_1, A_2, \dots, A_m and B_1, B_2, \dots, B_m be subsets of an n -element set, such that $A_i \cap B_j \neq \emptyset$ iff $i = j$. Then

$$m \cong \left(f \left(\frac{\sum_{i=1}^m |A_i|}{m}, \frac{\sum_{i=1}^m |B_i|}{m} \right) \right)^{-1} \cong \left(f \left(\max_{1 \leq i \leq m} |A_i|, \max_{1 \leq i \leq m} |B_i| \right) \right)^{-1}.$$

PROOF. Using Lemmas 2 and 1, we obtain

$$m \cdot f \left(\max_{1 \leq i \leq m} |A_i|, \max_{1 \leq i \leq m} |B_i| \right) \cong m f \left(\frac{\sum_{i=1}^m |A_i|}{m}, \frac{\sum_{i=1}^m |B_i|}{m} \right) \cong \sum_{i=1}^m f(|A_i|, |B_i|) \cong 1,$$

which gives our statement.

LEMMA 3. Let U_1, U_2, \dots, U_n and V_1, V_2, \dots, V_n be non-empty subsets of $Y = \{y_1, y_2, \dots, y_m\}$ such that $P_1 = \bigcup_{k=1}^n (U_k \times V_k)$ holds, where

$$P_1 = \{p : p = (y_i, y_j), i \neq j, 1 \leq i, j \leq m\}$$

Then for the unique λ , defined by $m = 1/f\left(\frac{\lambda}{2}, \frac{\lambda}{2}\right) : \sum_{k=1}^n (|U_k| + |V_k|) \cong \lambda \cdot m$ holds, where $f(x, y)$ is defined in Lemma 2.

PROOF. The following two statements are equivalent

- (i) $\bigcup_{k=1}^n (U_k \times V_k) = P_1 = \{p : p = (y_i, y_j), i \neq j, 1 \leq i, j \leq m\}$.
 (ii) The sets $A_i = \{k : y_i \in U_k\}$ and $B_i = \{k : y_i \in V_k\}$ ($1 \leq i \leq m$) have the property, that $A_i \cap B_j \neq \emptyset$ iff $i = j$.

(i \rightarrow ii) If $i \neq j$, $1 \leq i, j \leq m$ then $p = (y_i, y_j)$ is an element of $P_1 = \bigcup_{k=1}^n (U_k \times V_k)$, that is, there exists a k , $1 \leq k \leq n$ for which $y_i \in U_k$ and $y_j \in V_k$. Thus $k \in A_i \cap B_j$ and so $A_i \cap B_j \neq \emptyset$. On the other hand, if $A_i \cap B_j \neq \emptyset$, then there exists a k , $1 \leq k \leq n$ for which $k \in A_i \cap B_j$ that is $p = (y_i, y_j)$ is an element of $(U_k \times V_k)$ which means that $p = (y_i, y_j) \in P_1$ and so $i = j$.

(ii \rightarrow i) Assume $p = (y_i, y_j) \in P_1$, for some $1 \leq i, j \leq m, i \neq j$ follows. Thus, there exists a k , $1 \leq k \leq n$ for which $k \in A_i \cap B_j$. Then $p = (y_i, y_j)$ is an element of $(U_k \times V_k)$ which means that $p \in \bigcup_{k=1}^n (U_k \times V_k)$.

If $p = (y_i, y_j) \in \bigcup_{k=1}^n (U_k \times V_k)$ then there exists a k , $1 \leq k \leq n$ for which (y_i, y_j) is an element of $(U_k \times V_k)$. Thus $k \in A_i \cap B_j \neq \emptyset$ and so $i = j$, that is, $p \in P_1$.

Let us now suppose that $\sum_{k=1}^n (|U_k| + |V_k|) < \lambda m$ holds. Since $\sum_{k=1}^n (|U_k| + |V_k|) = \sum_{i=1}^m (|A_i| + |B_i|)$ is true, using the notation

$$\mu = \frac{\sum_{k=1}^n (|U_k| + |V_k|)}{m},$$

$$m = 1/f\left(\frac{\lambda}{2}, \frac{\lambda}{2}\right) > 1/f\left(\frac{\mu}{2}, \frac{\mu}{2}\right) \cong \left(f\left(\frac{\sum_{i=1}^m |A_i|}{m}, \frac{\sum_{i=1}^m |B_i|}{m}\right) \right)^{-1} \cong m$$

hold because on the base of Lemma 2 $f(x, y)$ is decreasing in both x and y , just as $f(x, y)$ is convex and symmetric ($f(x, y) = f(y, x)$) and so $f\left(\frac{x+y}{2}, \frac{x+y}{2}\right) \cong f(x, y)$. Finally the last inequality follows from Corollary 2. This is a contradiction. The proof is completed.

THEOREM 1. Put $Y=Z$, $|Y|=|Z| = \binom{l}{2} = m$ and $P_1 = \{p: p = (y_i, y_j), i \neq j, 1 \leq i, j \leq m\}$. Then the complexity of $P_1: \pi(P_1) = l \cdot \binom{l}{2}$.

PROOF. From Lemma 3 it follows $\pi(P_1) \geq \lambda m$. This lower bound is the best possible if $\lambda = l$ is an integer. Namely let A_i 's be all the $\binom{l}{2}$ element subsets of the set $X = \{1, 2, \dots, l\}$ and $B_i = A_i^c = X \setminus A_i$. Thus, for A_i 's and B_i 's (ii) of Lemma 3 holds. In this way for

$$U_k = \{y_i: k \in A_i\}, \quad V_k = \{y_i: k \in B_i\}$$

(i) $P_1 = \bigcup_{k=1}^l (U_k \times V_k)$ also holds. Accordingly if $\lambda = l$ then

$$\begin{aligned} \lambda \cdot m &= l \cdot \binom{l}{2} = l/f\left(\binom{l}{2}, \binom{l}{2}\right) = l/f\left(\frac{l}{2}, \frac{l}{2}\right) = \sum_{i=1}^m (|A_i| + |B_i|) = \\ &= \sum_{k=1}^l (|U_k| + |V_k|) \cong \pi(P_1) \cong \lambda \cdot m, \quad \text{and so} \quad \pi(P_1) = \lambda \cdot m. \end{aligned}$$

The proof is completed

2. Complexity of isosceles right triangles

In case of $Y=Z$ and $P_2 = \{p: p = (y_i, y_j), i < j, 1 \leq i, j \leq m\}$ let us compute the value of $\pi(P_2)$, defined in our introduction.

First I shall prove a lemma which uses the idea of Th. 1. in [3] and I shall draw two conclusions of it.

LEMMA 4. Let A_1, A_2, \dots, A_m and B_1, B_2, \dots, B_m be subsets of an n -element set X , such that $A_i \cap B_j \neq \emptyset$ iff $i < j$ ($1 \leq i, j \leq m$).

Then $\sum_{i=1}^m 2^{-(|A_i|+|B_i|)} \leq 1$ holds.

PROOF. For arbitrary pair $i < j$, ($1 \leq i, j \leq m$) $A_i \cap B_j \neq \emptyset$ holds. Then there is no $A \subset X$, for which both $A_i \subset A$, $A \cap B_i = \emptyset$ and $A_j \subset A$, $A \cap B_j = \emptyset$ hold, because otherwise $\emptyset \neq A_i \cap B_j \subset A \cap B_j = \emptyset$ would be true. Since there are $2^{(n-|A_i|-|B_i|)}$ sets A satisfying $A_i \subset A$ and $A \cap B_i = \emptyset$, thus $\sum_{i=1}^m 2^{(n-|A_i|-|B_i|)} \leq 2^n$ holds which gives our lemma.

LEMMA 5. Let A_1, A_2, \dots, A_m and B_1, B_2, \dots, B_m be subsets of an n -element set X such that $A_i \cap B_j \neq \emptyset$ iff $i < j$ ($1 \leq i, j \leq m$).

Then $m \cdot \log_2 m \leq \sum_{i=1}^m (|A_i| + |B_i|)$.

PROOF: Since 2^{-x} is a convex function and by Lemma 4

$$m \cdot 2^{-\sum_{i=1}^m \frac{(|A_i|+|B_i|)}{m}} \leq \sum_{i=1}^m 2^{-(|A_i|+|B_i|)} \leq 1 \text{ hold,}$$

which proves our lemma.

LEMMA 6. Let U_1, U_2, \dots, U_n and V_1, V_2, \dots, V_n be non-empty subsets of $Y = \{y_1, y_2, \dots, y_m\}$ such that

$$\bigcup_{k=1}^n (U_k \times V_k) = P_2 = \{p: p = (y_i, y_j), i < j, 1 \leq i, j \leq m\}.$$

Then $\sum_{k=1}^m (|U_k| + |V_k|) \geq m \cdot \log_2 m$.

PROOF. The following two statements are equivalent:

- (i) $\bigcup_{k=1}^n (U_k \times V_k) = P_2 = \{p: p = (y_i, y_j), i < j, 1 \leq i, j \leq m\}$,
- (ii) The sets $A_i = \{k: y_i \in U_k\}$ and $B_i = \{k: y_i \in V_k\}$ ($1 \leq i \leq m$) have the property, that $A_i \cap B_j \neq \emptyset$ iff $i < j$.

The proof of this equivalence is the same as it was in Lemma 3, the only difference is, that instead of $i \neq j$ you have to think of $i < j$.

Since $\sum_{k=1}^m (|U_k| + |V_k|) = \sum_{i=1}^m (|A_i| + |B_i|)$ is true, Lemma 5 gives our Lemma.

THEOREM 2. Put $Y=Z$, $|Y|=|Z|=2^l=m$ and $P_2=\{p:p=(y_i, y_j), i<j, 1\leq i, j\leq m\}$. Then the complexity of $P_2:\pi(P_2)=l\cdot 2^l$.

PROOF. From Lemma 6 it follows $\pi(P_2)\cong m\cdot \log_2 m$. This lower bound is the best possible if $\log_2 m=l$ is an integer. Namely let $Y=\{0, 1\}^l$ and $E=\bigcup_{i=0}^{l-1} \{0, 1\}^i$ be. If $\varepsilon=(\varepsilon_1, \varepsilon_2, \dots, \varepsilon_l)\in E$ let

$$W_{\varepsilon_0} = \{\delta\in\{0, 1\}^l: \delta = (\delta_1, \delta_2, \dots, \delta_l), \delta_j = \varepsilon_j \quad (1 \leq j \leq l), \delta_{l+1} = 0\}$$

$$W_{\varepsilon_1} = \{\delta\in\{0, 1\}^l: \delta = (\delta_1, \delta_2, \dots, \delta_l), \delta_j = \varepsilon_j \quad (1 \leq j \leq l), \delta_{l+1} = 1\}$$
 be.

In this case $P_2 = \bigcup_{\varepsilon\in E} (W_{\varepsilon_0}\times W_{\varepsilon_1})$ and $\sum_{\varepsilon\in E} |W_{\varepsilon_0}| + |W_{\varepsilon_1}| = l\cdot 2^l$ trivially hold. The proof is completed.

3. Complexity of Hadamard-matrices

Assume now, that $Y=Z=\{0, 1\}^l$ and $P_3=\{p:p=(y_i, y_j), y_i*y_j=0\}$ (where $y_i*y_j\equiv \sum_{k=1}^l y_{ik}\cdot y_{jk} \pmod{2}$). Let us compute the value $\pi(P_3)$, defined in our introduction. Let us remark that P_3 's are the Hadamard-matrices [1].

First I shall state a simple lemma and by means of it I shall give only a lower bound for $\pi(P_3)$. We introduce some notations:

If $U, V\subset Y$, $U*V=\{u*v: u\in U, v\in V\}$.

If $W\subset Y$, let $[W]$ be the subspace spanned by W , and finally $\dim W=\dim [W]$.

LEMMA 7. Let U, V be subsets of the vector space $Y=\{0, 1\}^l$ satisfying $U*V=0$. Then $\dim U+\dim V\leq l$ holds.

PROOF. U has $\dim U$ linearly independent vectors $\{u_i\}_{i=1}^{\dim U}$. For arbitrary $x\in\{0, 1\}^l$ and $x*U=0$ the $x*u_i\equiv 0 \pmod{2}$ $1\leq i\leq \dim U$ system of equations holds, which has $(l-\dim U)$ linearly independent solutions by the Cramer rule. The proof is completed.

THEOREM 3. Put $Y=Z=\{0, 1\}^l$ and $P_3=\{p:p=(y_i, y_j), y_i*y_j\equiv 0 \pmod{2}\}$ (where $y_i*y_j\equiv \sum_{k=1}^l y_{ik}\cdot y_{jk} \pmod{2}$). Then the complexity of P_3 :

$$\pi(P_3) \cong (2^l+1)2^{\frac{l}{2}}.$$

PROOF. Let U_1, U_2, \dots, U_n and V_1, V_2, \dots, V_n be non-empty subsets of Y , such that $P_3 = \bigcup_{k=1}^n (U_k \times V_k)$. Then $\sum_{k=1}^n |U_k| \cdot |V_k| \cong |P_3|$. Thus the following inequality holds:

$$\begin{aligned} (1) \quad \sum_{k=1}^n (|U_k| + |V_k|) &= \sum_{k=1}^n (|U_k|^{-1} + |V_k|^{-1}) \cdot |U_k| \cdot |V_k| \cong \\ &\cong \min_{1\leq k\leq n} (|U_k|^{-1} + |V_k|^{-1}) \sum_{k=1}^n |U_k| \cdot |V_k| \cong \min_{1\leq k\leq n} (|U_k|^{-1} + |V_k|^{-1}) \cdot |P_3|. \end{aligned}$$

On the other hand.

$$(2) \quad |U_k|^{-1} + |V_k|^{-1} \cong 2^{-\dim U_k} + 2^{-\dim V_k} \cong 2^{-\dim U_k} + 2^{-l + \dim U_k} \cong 2^{\left(1 - \frac{l}{2}\right)}$$

follows from Lemma 7.

Let us prove now

$$(3) \quad |P_3| = 2^{l-1}(2^l + 1) \text{ by induction over } l.$$

If $l=1$, then $|P_3|=3$. Let us suppose that (3) is true for $l-1$. If y_{ii} and y_{jj} are given and $y_{ii} \cdot y_{jj} = 0$ then $\sum_{k=1}^l y_{ik} y_{jk} \equiv 0 \pmod{2}$ in $2^{l-2}(2^{l-1}+1)$ cases by our inductual assumption. If $y_{ii} \cdot y_{jj} = 1$ holds then $\sum_{k=1}^l y_{ik} y_{jk} \equiv 0 \pmod{2}$ in $2^{2(l-1)} - 2^{l-2}(2^{l-1}+1)$ cases. Since $y_{ii} y_{jj} = 0$ in 3 cases thus $|P_3| = 3 \cdot 2^{l-2}(2^{l-1}+1) + 2^{2(l-1)} - 2^{l-2}(2^{l-1}+1) = 2^{l-1}(2^l+1)$ for l . From the formulas (1)–(3)

$$\sum_{k=1}^n (|U_k| + |V_k|) \cong \min_{1 \leq k \leq n} (|U_k|^{-1} + |V_k|^{-1}) |P_3| \cong 2^{1 - \frac{l}{2}} 2^{l-1}(2^l + 1).$$

The proof is completed

4. Complexity of the most complex lattice-configurations

I shall formulate an asymptotic lower bound for $\pi(m)$ which is a simple consequence of LUPANOV's Th. 4 in [5].

$$\text{THEOREM 4. } \pi(2^l) \gtrsim \frac{2^{2l}}{l \cdot \log_2 l}, \text{ where } a_i \gtrsim b_i \text{ iff } \liminf_l \frac{a_l}{b_l} \cong 1.$$

Open problems

The problems of Theorems 3 and 4, which are only partially solved, remained open questions:

1. Let

$$Y = \{y_1, y_2, \dots, y_{2^l}\} = \{0, 1\}^l$$

and

$$P_3 = \{p: p = (y_i, y_j), 1 \leq i, j \leq 2^l, y_i * y_j \equiv 0 \pmod{2}\}$$

be where $y_i * y_j \equiv \sum_{k=1}^l y_{ik} \cdot y_{jk} \pmod{2}$ and let us give a system U_1, U_2, \dots, U_n ;

V_1, V_2, \dots, V_n for which $P = \bigcup_{k=1}^n (U_k \times V_k)$ and

$$\sum_{k=1}^n (|U_k| + |V_k|) = \text{minimal},$$

where $U_k \times V_k$ is the Cartesian product of U_k and V_k .

2. Let $Y = \{y_1, y_2, \dots, y_m\}$ be and let us denote by $\pi(m)$ the minimal number that for every $P \subset (Y \times Y)$ there exists a system $U_1, U_2, \dots, U_n; V_1, V_2, \dots, V_n$ such that

$$P = \bigcup_{k=1}^n (U_k \times V_k) \text{ and}$$

$$\pi(m) \cong \sum_{k=1}^n (|U_k| + |V_k|) \text{ hold,}$$

where $(U \times V)$ is the Cartesian product of U and V . What is the exact value of $\pi(m)$?
I am greatly indebted to G. O. H. KATONA for his helpful advices.

REFERENCES

- [1] BERLEKAMP, E. R.: *Algebraic Coding Theory*, McGraw Hill, New York, 1969, 316—317.
- [2] KATONA, G. O. H.: On a Conjecture of Ehrenfeucht and Mycielski, *J. Combinatorial Th.* (to appear).
- [3] KATONA, G. and SZEMRÉDI, E.: On a Problem of Graph Theory, *Studia Sci. Math. Hungar.* 2 (1967), 23—28.
- [4] LUBELL, D.: A Short Proof of Sperner's Lemma, *J. Combinatorial Th.* 1 (1966) 299.
- [5] LUPANOV, O. B.: On Some Classes of Control Systems, "*Problemi Kybernetiki*" M., *Fizmat* 2 (1962) 63—97 (in Russian).
- [6] MESHALKIN, L. D.: A Generalization of Sperner's Theorem on the Number of Subsets of a Finite Set, *Teor. Verоятnost. i Primenen.* 8 (1963) 219—220 (in Russian).
- [7] TARJÁN, T. G.: On Complexity of Switching Circuits, *Problems of Control and Information Theory*, 3 (1974) 183—196.
- [8] YAMAMOTO, K.: Logarithmic Order of Free Distributive Lattices, *J. Math. Soc. Japan*, 6 (1954) 343—353.

MTA Közgazdaságtud. Intézet (Institute of Economics, Hungarian Academy of Sciences), 1051
Budapest, Münnich F. u. 7.

(Received February 22, 1975)