The problem

 \mathcal{A} : a class of labelled graphs, closed under isomorphism.

• What is the size of A_n = the set of graphs in \mathcal{A} on the vertex set $\{1, \ldots, n\}$?

► What is the structure of a typical graph in \mathcal{A}_n ?

Note: in general, lots of dependence in the edges of a graph drawn uniformly at random from $\mathcal{A}_n!$

 $\mathcal{B} =$ fixed finite set of graphs. A graph is bad, if it has a minor (\nearrow) in \mathcal{B} . Ex \mathcal{B} = graphs that are not bad (no minor in \mathcal{B}). We study classes $E_{x}(k+1)\mathcal{B}$ of graphs that do not have k+1vertex-disjoint bad subgraphs (k fixed).

Example. Ex(k + 1){ K_3 } is the class of all graphs that do not have k + 1 vertex disjoint cycles. This and similar classes were studied before (\downarrow) .

Classes in the current work include the case $\mathcal{B} = \{K_4\}$, the graphs which do not contain

$$\boxtimes \boxtimes \ldots \boxtimes k+1$$
 times

as a minor.

Small blockers

 \mathcal{B} -blocker Q of G = every bad subgraph of G meets Q in at least one point = G - Q is not bad.

Fans = graphs consisting of a path and an apex vertex connected to each vertex on the path.



Theorem. (V.K. and C. McDiarmid, 2011). If the graphs in \mathcal{B} are 2connected and $\mathcal{A} = \operatorname{Ex} \mathcal{B}$ does not contain all fans, then all but an exponentially small proportion of graphs without k + 1 disjoint bad subgraphs have a *B*-blocker of size k.

The number of such graphs on the vertex set $\{1, \ldots, n\}$ is $c(k, \mathcal{B})2^{kn}|\mathcal{A}_n|\left(1-e^{-\Omega(n)}\right).$

Graph minors

H is a minor of G if it can be obtained from G by a series of edge contractions/deletions.



Minor-closed classes are subject to fundamental graph theory results. For example, Kuratowski's theorem states that planar graphs are exactly those graphs that do not have $K_{3,3}$ or K_5 as a minor. Robertson's and Seymour's theory of graph minors proves that every minor-closed class of graphs has a finite list of forbidden minors.

For the class of planar graphs the two main questions (\nwarrow) have been studied quite intensively in the past decade (Bender, Bernardi, Bodirsky, Drmota, Fusy, Gao, Gimenez, Kang, McDiarmid, Noy, Steger, Panagiotou, Welsh, Wormald...). For this class, and for classes with 2-connected excluded minors, the asymptotic growth and the properties of random graphs are now quite well understood.

In other words, the simplest possible type of structures overhelmingly dominates the class of all graphs without k+1 disjoint minors in \mathcal{B} .

Example. A random graph on vertex set $\{1, \ldots, n\}$ without k + 1 disjoint cycles is "almost" a uniformly random forest F along with k "apex" vertices, where the edges between the apex vertices are independent and exist with probability 1/2. A uniformly random graph without 3 disjoint cycles:



Growth constants

 ${\cal A}$ has a growth constant γ if

 $(|\mathcal{A}_n|/n!)^{1/n} \to \gamma.$

 $\overline{\gamma}(\mathcal{A})$ =upper limit of the above = $\rho(\mathcal{A})^{-1}$ = inverse radius of convergence of the EGF of \mathcal{A} .

Conjecture (Bernardi, Noy and Welsh, 2010). Every minor-closed class of graphs has a growth constant.

Redundant \mathcal{B} -blocker Q of G = every bad subgraph of G meets Q in at least two points. $rd_{I}B = the class of graphs that have a redundant$ \mathcal{B} -blocker of size at most I.

 $\operatorname{rd}_{2k+1}\mathcal{B} \subseteq \operatorname{Ex}(k+1)\mathcal{B}$: each bad subgraph in G uses at least two vertices from the blocker, so we can form no more than k disjoint such subgraphs.

Theorem 1. If $\operatorname{Ex} \mathcal{B}$ contains all fans and condition A holds then there is a constant $k_0 = k_0(\mathcal{B})$ such that for all $k \ge k_0$

$$\overline{\gamma}(\operatorname{Ex}(k+1)\mathcal{B}) = \overline{\gamma}(\operatorname{\mathsf{rd}}_{2k+1}\mathcal{B})$$

Examples of suitable \mathcal{B} : $\{K_{2,3}\}$, $\{K_4\}$, $\{W_5\}$ and $\{K_{3,3}, P_0\}$. $W_5 =$ wheel on 5 vertices; P_0 = arbitrary planar graph.

For \mathcal{B} and k as in Theorem 1, there are arbitrarily large n such that in almost all graphs on $\{1, \ldots, n\}$ which do not have k+1 bad subgraphs, we need at least 2k vertices to block all bad subgraphs. This is twice more than for \mathcal{B} considered earlier (\swarrow).

Theorem 2. If $\operatorname{Ex} \mathcal{B}$ contains all fans and conditions A and B hold then there is a constant $k_0 = k_0(\mathcal{B})$ such that for all $k \ge k_0$ the class $\operatorname{Ex}(k+1)\mathcal{B}$ has a growth constant.

Examples of suitable \mathcal{B} : { K_4 }, { W_5 }.

Condition A

The graphs in \mathcal{B} are 2-connected, there is at least one planar graph and there is a finite number m such that if we take any tree T of size *m* and add two apex vertices connected to each vertex of T, then we obtain a bad graph.

Condition B

There are finite numbers m and l such that if we take a 2-connected graph H of size m and a new vertex with I neighbours in V(H), then we obtain a bad graph.

Graphs with no k + 1 **disjoint minors** K_4

Theorem 3. Let $A = Ex(k+1)\{K_4\}$. For any $k = 1, 2, \cdots$

- \mathcal{A} has a growth constant γ_k .
- $\gamma_1 = \gamma(\text{Ex} \boxtimes \boxtimes) = 23.5241..$
- A uniformly random graph from A_n whp has a unique redundant blocker Q of size 2k + 1.
- ► Each vertex of *Q* whp has a linear degree.
- There is $c_k > 0$ such that $|A_n| = c_k n^{-5/2} n! \gamma_k^n (1 + o(1))$.

Proof elements

- ► Theory of graph minors.
- points, must meet S.
- New structural results
- Analytic combinatorics.
- by objects of different types.

Typical graphs without two disjoint minors K_4



Leaf-like shapes = arbitrary series-parallel *networks*, most of them small. Above is just a "core" of a graph in $Ex 2K_4$ (it is typically of linear size). Attach arbitrary rooted series-parallel graphs at arbitrary vertices of this; add a redundant blocker - $\{x, y, z\}$ - and connect x to each red point, y to each green point and z to each blue point. Unlike in planar graphs, there are vertices of different types: colours can occur only at the "joints"; only the joints of the grey blocks can have arbitrary colour. I.e., the top-right vertex cannot have colour blue, otherwise z alone would form a minor K_4 , and $\{x, y, z\}$ would not be redundant.

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Robertson and Seymour proved that any $G \in Ex(k+1)B$ has a constant size \mathcal{B} -blocker, as long as \mathcal{B} contains a planar graph.

• Earlier results on planar and other minor-closed classes.

Ideas of McDiarmid, Steger and Welsh (2005) help to *avoid* generating functions in the proof of Theorem 2. A new version of a lemma of V.K. and McDiarmid (2011): every $G \in Ex(k+1)B$ can be represented by (a) a constant-size set Q (b) a set $S \subseteq Q$ of size at most 2k and (c) a graph in $\operatorname{Ex} \mathcal{B}$, so that each bad sugraph that meets Q in at most two

If G - Q is connected, we can *split* (= reverse contraction) a constant number of vertices in G and represent G as O(1) graphs in rd_{2k+1} B and O(1) edges (the crucial point).

Easy by the "First principle" of Analytic Combinatorics: $Ex(2k+1)B \approx 1$ a class of sets of graphs composed of O(1) elements from $\operatorname{rd}_{2k+1}\mathcal{B} \implies$ $\rho(\mathrm{Ex}\,(k+1)\mathcal{B}) = \rho(\mathrm{rd}_{\,2k+1}\,\mathcal{B}).$

More analytic combinatorics

Bodirsky, Gimenez, Kang and Noy (2005) obtained expressions for the EGF of series-parallel graphs = $Ex \{K_4\}$. Using these, we analyse the structure (\downarrow) of graphs with small redundant blockers. This includes counting trees where leaves, internal vertices and edges are substituted